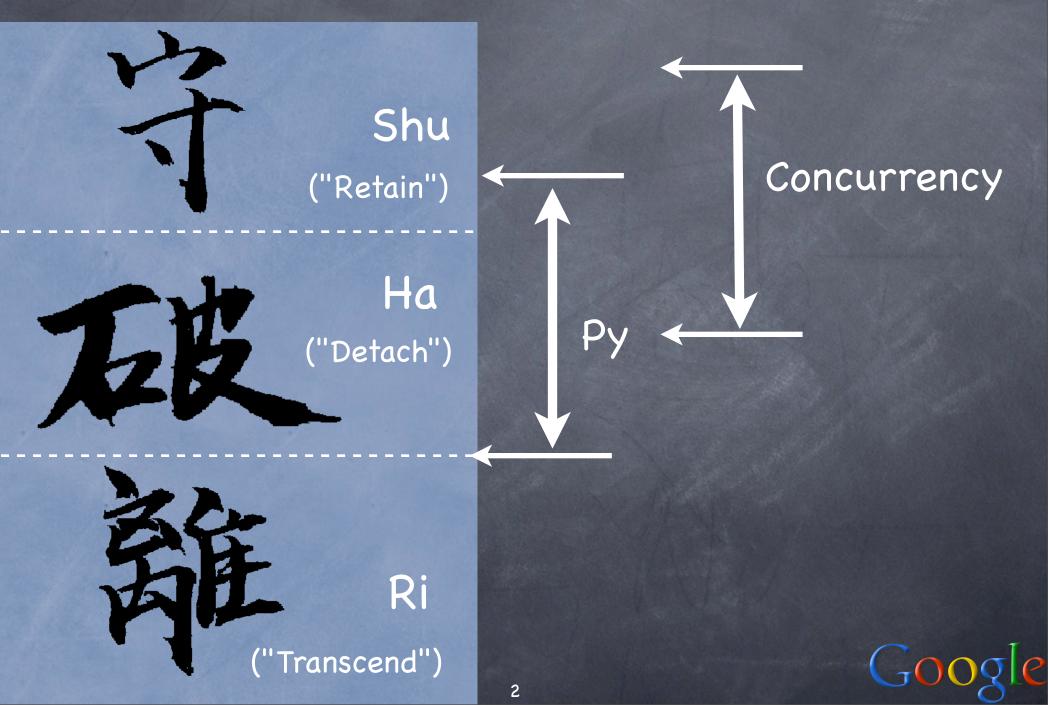
Python Patterns of Concurrent Programming

http://www.aleax.it/accu_pyconc.pdf

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The "levels" of this talk



Concurrency

The want to do several things "at once" to react to "events" (mostly, but not exclusively, "external" ones) @ actions by the user[s] (on a UI, or...) In network events, I/O completion, OS... ø parsing, graph-walking, ... To reduce "latency" (vs "throughput") to exploit many available resources (CPUs, cores, computers...) towards a single goal simultaneously "parallel", "distributed", ...

Event-driven: the callback

The "callback" idea Shades of the Observer DP ø pass a callable to "somebody" In Observer: to the Observable target The "somebody" stores it "somewhere" @ a container, an attribute, whatever and will call it "when appropriate" In Observer: on state-changes In Event-driven: on meaningful "events" also used in "customization" (e.g., sort)

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Callback issues

So what arguments are to be used on the call? o none: simplest, very rough In ODP: the Observable object whose state just changed Iets 1 callable observe several Obs'bles or: "description" of state changes saves "round-trips" to obtain them In EDA: identifier or description of event ø but -- what about other arguments (related to the callable, not to the Obs'ble/Event)...?

Fixed args in callbacks

- functools.partial(callable, *a, **kw)
 pre-bind any or all arguments
- ø however...:

 - @ x.set_cb(f, *a, **kw)

...having the set-callback itself accept (and pre-bind) arguments is far neater/handier
 sombunall¹ Python callback systems do that

6

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¹: Robert Anton Wilson

Callback "dispatching"

what if more than one callback is set for a single event (or, Observable)? remember and call the latest one only simplest, roughest remember and call them all ø how do you _remove_ a callback? can one callback preempt others? can events (or state changes) be "grouped"? use object w/methods instead of callable

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Callbacks and Errors

@ are "errors" events like any others? or are they best singled-out? http://www.python.org/pycon/papers/deferex/ The Deferred pattern: one Deferred holds... N "chained" callbacks for "successes" + M "chained" callbacks for "errors" each callback is held WITH opt *a, **kw
 In plus, argument for "event / error identification" (or, result of previous callback along the appropriate "chain")

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Scheduled callbacks

standard library module sched s = sched.Sched(timefunc, delayfunc) @ e.g, Sched(time.time, time.sleep) @ evt = s.enter(delay, priority, callable, arg) ø or s.enterabs(time, priority, callable, arg) may s.cancel(evt) later
 s.run() runs events until queue is empty (or an exception is raised in callable or delayfunc: it propagates but leaves s in stable state, s.run can be called again later)

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Some other Python callbacks

o for system-events: @ atexit.register(callable, *a, **k) Idhandler = signal.signal(signum, callable) sys.displayhook, sys.excepthook, sys.settrace(callable) @ readline.set_startup_hook, set_pre_input_hook, set_completer ø parsing, timing, debugging, customization, ...



Event-driven: "the Loop"

ø demultiplex external events coming into the program from several distinct channels select.select, select.poll (+ 3rd-party ones: python-epoll, Py-Kqueue, ...) win32all.MsgWaitForMultipleEvents "register" event sources (fd's, windows, kernel sync objects,) I call the demux function (maybe w/timeout) ø returns set of events occurred (or, times out, allowing optional further "polling"...)

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Dispatching Events

The "Event Loop" per se ``bottlenecks" events into 1 spot (+ others if ``modal") Solution: if/else tree at that spot ø better: register callbacks for "specific" events (by type, source, or ...) & add to the "Event Loop" the registry and dispatching of callbacks it affords Loop + Registry/Dispatching == REACTOR
 http://www.cs.wustl.edu/~schmidt/PDF/ reactor-siemens.pdf

asyncore

asyncore.loop(timeout, use_poll, map, count) only supports events handled by select
(or poll, if use_poll is true) also does dispatching (so: a reactor) items in map (or asyncore.socket_map) must have methods readable, writable, all of socket's methods, plus handle_* ones... In normally subclass asyncore.dispatcher also handles add-to-the-map, wraps socket methods To buffer: dispatcher_with_send ø for files (Unix only): file_dispatcher

asyncore Echo srv (1/2)

import asyncore

class MainServerSocket(asyncore.dispatcher): def __init__(self, port): asyncore.dispatcher.__init__(self) self.create_socket(socket.AF_INET, socket.SOCK_STREAM) self.bind(('',port)) self.listen(5) def handle_accept(self): newSocket, address = self.accept() SecondaryServerSocket(newSocket)

asyncore Echo srv (2/2)

MainServerSocket(8881)
asyncore.loop()



Twisted Core

Twisted.internet.reactor module IReactorCore: "system events", run, stop IReactorTCP: listen/connect w/Factory (ServerFactory, ClientFactory) Sectory it f has buildProtocol method IProtocol has dataReceived method many other reactor itf's: FDSet, Process,
 SSL, Threads, Time, UDP, UNIX many implementations: select, poll, KQueue, Win32 (WFMO/IOCP) + many with GUI event-loops integration

Twisted Echo server

from twisted.internet import protocol, reactor

class EchoProtocol(protocol.Protocol):
 def dataReceived(self, data):
 self.transport.write(data)

factory = protocol.Factory()
factory.protocol = EchoProtocol

reactor.listenTCP(8881, factory)
reactor.run()



Interleaving execution

Traditional sequential processing: $data = get_all_data()$ results = process_all_data(data) show_all_results(results) finished = False while not finished: d, $f = get_some_data()$ r, f = process_some_data(d, f) finished = show_some_results(r, f) reduces latency (may worsen throughput)
 AKA "µthreads", "fibers", "tasklets"...

Interleaving issues

ø each component must: restore its previous state (if any), ø do "some" work (not TOO much!), save its current state (if not finished), ø yield control to "other components" In note: MUST yield if risks blocking (I/O...) Icear pluses: programmer has complete control, works great w/event-driven approaches, no "surprises", ... issues: lots of work, some "delicate" points

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Python interleaving

may use classes (to group state & behavior) generators are very handy for this automatically "preserve state" including "point of execution" at yield time ø yield is now "bidirectional" (expression) In the standard "scheduling" conventions (yet) but, cfr <u>http://mail.python.org/pipermail/</u>
 python_list/2007_May/438370.html
 ø don't confuse microthreading, coroutines, and continuations...

Stackless Python

stasklets (AKA microthreads): wrap a function for interleaving purposes supply critical sections, scheduling (cooperative _or_ even preemptive...) channels: let tasklets send/receive data Intrinsically cooperate w/scheduling serialization (pickling/unpickling) for checkpointing, transferring tasks, ... ø popular in games, online and non (EVE) Online, Mythos, Sylphis 3D, ...) for speed and convenience

"Real" Threads

The threading module (NOT legacy thread) Substitution of the second state of the sec threading facilities for Jython, IronPython) OPython adds a Global Interpreter Lock
 OPython adds
 OPython
 (GIL) to ease integrating non-threadaware/safe external C libraries C-coded extensions may explicitly release the GIL ("allow threading")/re-acquire it semantics constrained by cross-platform

needs (priorities, thread-interruption...)

Threads & "Atomicity" # will this work? why, or, why not? d = {}; fd = open('fil.txt') def f(): for L in fd: k, v = L.split()[:2]d[k] = vt1 = threading.Thread(target=f)t2 = threading.Thread(target=f) t1.start(); t2.start() t1.join(); t2.join()

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"Thread-safe" iterator def tsiter(it, L=None): if L is None: L = threading.Lock()it = iter(it) while True: with L: yield it.next()

NB: in 2.5, this needs adding a: # from __future__ import with_statement

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"Thread-safe" mapping class tsdict(dict): def __init__(self, *a, **k): self.L = threading.Lock() with self.L: dict.__init__(self, *a, **k) def __setitem__(self, k, v): with self.L: dict.__setitem__(self, k, v) def __getitem__(self, k): with self.L: return dict.__getitem__(self, k)

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Big risks with this...

ø not covering _all_ "atomicity" needs: @ e.g.: "for k in tsiter(tsd): <body>" STILL needs "NO ALTERATION" to tsd throughout the loop (what ensures this?) The more, finer-grained locks are around, the higher the risk of deadlocks: T1 gets lock A then waits on lock B, T2 gets lock B then waits on lock A... orace conditions and deadlocks are the worst kinds of bugs (hard to reproduce, hard to test for, very hard to debug, ...)

A more structured approach

make every "shared resource" either UNCHANGING during multitasking, @ or, have it OWNED by ONE dedicated thread (only one changing OR accessing it) every other thread requests operations on the shared resource by sending MESSAGES to the dedicated owner-thread (and may wait for a result-message if applicable) Queue.Queue is the intrinsically-threadsafe communication structure for messages from thread to thread (work-request, result)

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TS file-read-to-Queue

a "self-activated" owner-thread variant

lines = Queue.Queue(N)def tsread(fn, linequeue): with open(fn) as fd: for line in fd: linequeue.put(line) t = threading.Thread(target=tsread,args=('fil.txt', lines)) t.setDaemon(True) t.start()



TS "dict service" (1/2)# a more classic "dedicated work thread" def tsmapping(wrq, d): while True: op, k, v, rq = wrq.get() if op == 'set': d[k] = velif op=='keys': v = d.keys() elif op=='in': v = k in d else: v = d.get(k, v)if rq: rq.put(v)



TS "dict service" (2/2)

...and a little syntax sugar on top...: class tsdict(object, UserDict.DictMixin): def __init__(self, *a, **k): self.wrg = Queue.Queue(N) t = threading.Thread(target=tsmapping, args=(self.wrq, dict(*a, **k))) t.setDaemon(True); t.start() self.rsq = Queue.Queue() def keys(self): self.wrq.put(('keys','','',self.rsq)) return self.rsq.get() ..._getitem__, __setitem__,... Goog

Granularity & Performance

@ each context-switch among threads has a performance cost (potentially in both latency and throughput) so does each locking (and Queue.Queue intrinsically does locking, too) consider "batching things up" a bit compromise betw. throughput & latency ø ideal points for context switches: where they would occur anyway (syscalls, I/O) avoid polling; consider thread-pools; mix and match threads w/event-driven ops; ...

Processes

Meavier/costlier than threads (but not by much, in Linux, Solaris, Mac OS X, BSD, ...) ø better isolation is a good guard vs bugs (and, can "drop privileges" &c for security) In o GIL -- use all CPUs/cores implicitly "resource sharing" goes a bit against the grain (possible, but "message-style" IPC mechanisms are generally preferable) IPC via sockets affords nearly unbounded potential scalability... see http://pypi.python.org/pypi/processing

"Orchestrating" concurrency

Twisted is good at orchestrating all this... ø focus on event-driven operations strong support for threads & processes ø particularly strong at networking some support for interleaving ("Flow" now deprecated, use task.Cooperator class) A highly structured (chiefly via interfaces) Iots of "moving parts" (alternatives, details, docs, abstractions...) Is there something simpler to use...?

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NetworkSpaces

Icose kin to Linda / tuple spaces Ø Python implementation uses Twisted Iclients also available for R (and Matlab) @ dual-licensed open source (GPL or Pro) 5 primitives: store, fetch{Try}, find{Try} store/fetch like Queue's put/get multiple "slots" w/arbitrary name per ws "find" for non-removing read access works with Sleigh (to distribute and load-balance work across cores/CPUs/nodes)

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nws basics

ø you need to be running an nws server There's only one kind (Python+Twisted) may run as a daemon may monitor it through a web itf
 ø optional, helpful "pybabelfish" daemon must have nws clients installed on all hosts may be Python, R, or Matlab ones Ianguage interop via ASCII strings, only ø within 1 language, any serializable obj OK

An NWS "server"

>>> from nws import client >>> ws = client.NetWorkSpace('primes') >>> import gmpy >>> pr = [gmpy.mpz(2)]>>> def store_prime(n): while len(primes)<=n:</pre> ••• pr.append(pr[-1].next_prime()) • • • ws.store('prime', int(primes[n]))

>>> while True:
 store_prime(ws.fetch('n'))

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An NWS "client"

>>> from nws import client
>>> ws = client.NetWorkSpace('primes')
>>> def getprime(n):
... pr.store('n', n)
... return pr.fetch('prime')

>>> print [getprime(n) for n in range(23)]
[2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31,
37, 41, 43, 47, 53, 59, 61, 67, 71, 73,
79, 83]

Sleigh

an enhancement of the original Simple Network Of Workstation (SNOW) idea a nws.sleigh.Sleigh instance s coordinates a set of "workers" may be local, or on arbitrary nodes @ uses ssh (or web, or ...) to start nodes may send the same task to each worker
 (s.eachWorker), and/or, may shard tasks among them (s.eachElem) @ also: s.imap, s.starmap, ... [use SleighArgs, *not* "SleighArguments"]

Parallel Python

similar to Sleigh (but more direct, less fancy iteration): message-based, supports both SMP and clusters semi-transparently Is class pp.Server supports ncpus, remote nodes, stats/logging, and: submit(func, args, depfuncs, modules, callback, callbackargs, group, globals) returns callable that wait & returns results also, for explicit use: wait(group)





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40